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Team automata : a formal approach to the modeling of collaboration between system components

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List of Symbols

2. Preliminaries

\subseteq	set inclusion, 23
\subset	proper set inclusion, 23
\setminus	set difference, 23
$\#$	cardinality (of a set), 23
\emptyset	the empty set, 23
$[n]$	shorthand for $\{1, 2, \dots, n\}$, 23
\mathbb{N}	set of positive integers, 23
Π	cartesian product (prefix notation), 23
\times	cartesian product (infix notation), 23
proj_j	projection on element j , 23
proj_J	projection on subset J , 23
$\text{proj}_j^{[2]}$	shorthand for $\text{proj}_j \times \text{proj}_j$, 24
$\text{proj}_J^{[2]}$	shorthand for $\text{proj}_J \times \text{proj}_J$, 24
$f \upharpoonright C$	restriction of function f to a subset C of its domain, 24
Σ	alphabet, 24
λ	the empty word, 24
$ w $	length (of a word w), 24
$w(i)$	i -th letter (of a word w), 24
$\#_a(w)$	total number of occurrences of letter a (in a word w), 24
$\text{alph}(w)$	alphabet (of a word w), 25
Σ^*	set of all finite words over Σ , 25
Σ^+	set of all nonempty finite words over Σ , 25
Σ^ω	set of all infinite words over Σ , 25
Σ^∞	set of all words over Σ , 25
$u \cdot v$	concatenation (of words u and v), 25
$K \cdot L$	concatenation (of languages K and L), 25
$\text{pref}(w)$	set of prefixes (of a word w), 26
$w[n]$	prefix of length n (of a word w), 25
$\lim_{n \rightarrow \infty} v_n$	limit (of words $v_1 \leq v_2 \leq \dots$), 26
pres_Γ	function preserving the symbols from Γ (and erasing all other symbols), 27

3. Automata

\mathcal{A}	automaton, 29
Q	set of states (of \mathcal{A}), 29
Σ	set of actions or alphabet (of \mathcal{A}), 29
δ	set of labeled transitions (of \mathcal{A}), 29
I	set of initial states (of \mathcal{A}), 29
δ_a	set of a -transitions (of \mathcal{A}), 30
$\mathbf{C}_{\mathcal{A}}$	set of finite computations of \mathcal{A} , 30
$\mathbf{C}_{\mathcal{A}}^\omega$	set of infinite computations of \mathcal{A} , 30
$\mathbf{C}_{\mathcal{A}}^\infty$	set of computations of \mathcal{A} , 30
$\mathbf{B}_{\mathcal{A}}^{\Theta, \infty}$	Θ -behavior of \mathcal{A} , 31
$\mathbf{B}_{\mathcal{A}}^\Theta$	finitary Θ -behavior of \mathcal{A} , 31
$\mathbf{B}_{\mathcal{A}}^{\Theta, \omega}$	infinitary Θ -behavior of \mathcal{A} , 31
Q_S	set of reachable states (of \mathcal{A}), 36
Σ_A	set of active actions (of \mathcal{A}), 36
δ_T	set of useful transitions (of \mathcal{A}), 36
$\mathcal{A}_1 \sqsubseteq \mathcal{A}_2$	containment (of \mathcal{A}_1 in \mathcal{A}_2), 36
\mathcal{A}_A^Θ	Θ -action-reduced version of \mathcal{A} , 37
\mathcal{A}_T^Θ	Θ -transition-reduced version of \mathcal{A} , 38
\mathcal{A}_S	state-reduced version of \mathcal{A} , 46
\mathcal{A}_A	action-reduced version of \mathcal{A} , 50
\mathcal{A}_T	transition-reduced version of \mathcal{A} , 50
\mathcal{A}_R	reduced version of \mathcal{A} , 50

4. Synchronized Automata

\mathcal{I}	index set, 59
\mathcal{A}_i	automaton, 59
\mathcal{S}	set of automata, 59
$\Delta_a(\mathcal{S})$	complete transition space of a in \mathcal{S} , 60
\mathcal{T}	synchronized automaton, 60
$SUB_J(\mathcal{T})$	the subautomaton of \mathcal{T} determined by J , 64
SUB_J	the subautomaton (of \mathcal{T}) determined by J , 64
$\pi_{\mathcal{A}_j}$	projection on automaton \mathcal{A}_j , 70
π_{SUB_J}	projection on subautomaton SUB_J , 70
\mathcal{D}	indexed set, 76
$\mathcal{V}(\mathcal{D})$	all finitely nested cartesian products of sets from \mathcal{D} , 76
$\text{dom}(V)$	domain of an element V , 76
u_V	function unpacking elements v from V , 77

$\langle v \rangle_V$	reordering of an element $v \in V$ relative to the construction of V , 77
$\langle \langle \mathcal{T} \rangle \rangle_{\mathcal{S}}$	reordered version of synchronized automaton \mathcal{T} (w.r.t. \mathcal{S}), 81
\mathcal{T}	synchronized automaton, 84
$Free(\mathcal{T})$	set of <i>free</i> actions of \mathcal{T} , 85
$AI(\mathcal{T})$	set of <i>ai</i> actions of \mathcal{T} , 85
$SI(\mathcal{T})$	set of <i>si</i> actions of \mathcal{T} , 86
$\mathcal{R}_a^{no}(\mathcal{S})$	predicate <i>no-constraints</i> , 88
$\mathcal{R}_a^{free}(\mathcal{S})$	predicate <i>is-free</i> for a in \mathcal{S} , 88
$\mathcal{R}_a^{ai}(\mathcal{S})$	predicate <i>is-ai</i> for a in \mathcal{S} , 89
$\mathcal{R}_a^{si}(\mathcal{S})$	predicate <i>is-si</i> for a in \mathcal{S} , 89
j	element of \mathcal{I} , 90
J	subset of \mathcal{I} , 90
Θ	arbitrary alphabet disjoint from set Q of states (of \mathcal{T}), 90

5. Team Automata

\mathcal{C}	component automaton, 116
Σ_{inp}	set of input actions or input alphabet (of \mathcal{C}), 116
Σ_{out}	set of output actions or output alphabet (of \mathcal{C}), 116
Σ_{int}	set of internal actions or internal alphabet (of \mathcal{C}), 116
$und(\mathcal{C})$	underlying automaton of \mathcal{C} , 116
Σ	set of actions or (full) alphabet (of \mathcal{C}), 116
Σ_{ext}	set of external actions or external alphabet (of \mathcal{C}), 116
Σ_{loc}	set of locally-controlled actions or locally-controlled alphabet (of \mathcal{C}), 117
$\mathbf{B}_{\mathcal{C}}^{\Sigma_{inp}, \infty}$	input behavior (of \mathcal{C}), 117
$\mathbf{B}_{\mathcal{C}}^{\Sigma_{out}, \infty}$	output behavior (of \mathcal{C}), 117
$\mathbf{B}_{\mathcal{C}}^{\Sigma_{int}, \infty}$	internal behavior (of \mathcal{C}), 117
$\mathbf{B}_{\mathcal{C}}^{\Sigma_{ext}, \infty}$	external behavior (of \mathcal{C}), 117
$\mathbf{B}_{\mathcal{C}}^{\Sigma_{loc}, \infty}$	locally-controlled behavior (of \mathcal{C}), 117
\mathcal{I}	index set, 118
\mathcal{C}_i	component automaton, 118
Σ_i	set of actions (of \mathcal{C}_i), 118
\mathcal{S}	set of component automata, 118
\mathcal{S}	composable system, 118
\mathcal{T}	team automaton, 120
$und(\mathcal{T})$	underlying synchronized automaton of \mathcal{T} , 120
$SUB_J(\mathcal{T})$	the subteam of \mathcal{T} determined by J , 122
SUB_J	the subteam (of \mathcal{T}) determined by J , 122

\mathcal{S}	composable system, 123
$\langle\langle \mathcal{T} \rangle\rangle_{\mathcal{S}}$	reordered version of team automaton \mathcal{T} w.r.t. \mathcal{S} , 125
\mathcal{T}	team automaton, 126
Σ_{inp}	set of input actions (of \mathcal{T}), 126
Σ_{out}	set of output actions (of \mathcal{T}), 126
Σ_{int}	set of internal actions (of \mathcal{T}), 126
Σ	set of actions (of \mathcal{T}), 126
Σ_{ext}	set of external actions (of \mathcal{T}), 126
Σ_{loc}	set of locally-controlled actions (of \mathcal{T}), 126
$\mathcal{I}_{a,inp}(\mathcal{S})$	input domain of a in \mathcal{S} , 126
$\mathcal{I}_{a,out}(\mathcal{S})$	output domain of a in \mathcal{S} , 126
$\mathcal{I}_{a,inp}$	input domain of a (in \mathcal{S}), 127
$\mathcal{I}_{a,out}$	output domain of a (in \mathcal{S}), 127
$SUB_{a,inp}(\mathcal{T})$	input subteam of a in \mathcal{T} , 127
$SUB_{a,out}(\mathcal{T})$	output subteam of a in \mathcal{T} , 127
$SUB_{a,inp}$	input subteam of a (in \mathcal{T}), 127
$SUB_{a,out}$	output subteam of a (in \mathcal{T}), 127
$SIPP(\mathcal{T})$	set of <i>sipp</i> actions of \mathcal{T} , 129
$WIPP(\mathcal{T})$	set of <i>wipp</i> actions of \mathcal{T} , 129
$SOPP(\mathcal{T})$	set of <i>sopp</i> actions of \mathcal{T} , 129
$WOPP(\mathcal{T})$	set of <i>wopp</i> actions of \mathcal{T} , 129
$MS(\mathcal{T})$	set of <i>ms</i> actions of \mathcal{T} , 131
$SMS(\mathcal{T})$	set of <i>sms</i> actions of \mathcal{T} , 131
$WMS(\mathcal{T})$	set of <i>wms</i> actions of \mathcal{T} , 132
$\mathcal{I}_{a,inp}$	input domain of a (in \mathcal{S}), 141
$\mathcal{I}_{a,out}$	output domain of a (in \mathcal{S}), 141
$\mathcal{R}_a^{sipp}(\mathcal{S})$	predicate <i>is-sipp</i> for a in \mathcal{S} , 141
$\mathcal{R}_a^{wipp}(\mathcal{S})$	predicate <i>is-wipp</i> for a in \mathcal{S} , 141
$\mathcal{R}_a^{sopp}(\mathcal{S})$	predicate <i>is-sopp</i> for a in \mathcal{S} , 142
$\mathcal{R}_a^{wopp}(\mathcal{S})$	predicate <i>is-wopp</i> for a in \mathcal{S} , 142
$\mathcal{R}_a^{ms}(\mathcal{S})$	predicate <i>is-ms</i> for a in \mathcal{S} , 144
$\mathcal{R}_a^{sms}(\mathcal{S})$	predicate <i>is-sms</i> for a in \mathcal{S} , 144
$\mathcal{R}_a^{wms}(\mathcal{S})$	predicate <i>is-wms</i> for a in \mathcal{S} , 144
$\Sigma_{i,ext}$	set of external actions (of \mathcal{C}_i), 150
$\Sigma_{i,loc}$	set of locally-controlled actions (of \mathcal{C}_i), 150
j	element of \mathcal{I} , 150
J	subset of \mathcal{I} , 150
$\Sigma_{J,ext}$	set of external actions (of SUB_J), 150
$\Sigma_{J,loc}$	set of locally-controlled actions (of SUB_J), 150

6. Behavior of Team Automata

pREG	family of prefix-closed regular finitary languages, 164
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REG	family of regular languages, 164
FIN	family of finite languages, 164
CA	$\{\mathbf{B}_C^\Sigma \mid C \text{ is a finite component automaton with alphabet } \Sigma\}$, 164
\mathbf{CA}^{alph}	$\{\mathbf{B}_C^{alph} \mid C \text{ is a finite component automaton}\}$ (with $alph \in \{inp, out, int, ext, loc\}$), 165
\mathcal{I}	index set, 166
C_i	component automaton, 166
Σ_i	set of actions (of C_i), 166
\mathcal{S}	composable system, 166
\mathcal{T}	team automaton, 166
Σ	set of actions (of \mathcal{T}), 166
Θ	arbitrary alphabet disjoint from set Q of states (of \mathcal{T}), 166
j	element of \mathcal{I} , 166
$uAI_j(\mathcal{T})$	set of useful j -ai actions (of \mathcal{T}), 169
\parallel	shuffle, 183
$\parallel\parallel$	fair shuffle, 183
$\ d\ $	norm (of decomposition d), 198
$\ \cdot\ _{i \in [n]}$	n -ary fair shuffle, 205
$\ \cdot\ _{i \in [n]}$	n -ary shuffle, 205
$\ \cdot\ ^\Gamma$	S-shuffle on Γ , 207
$\ \cdot\ ^\Gamma$	fair S-shuffle on Γ , 207
$alph(L)$	alphabet (of a language L), 208
$\Sigma_1 \underline{\parallel}_{\Sigma_2}$	fS-shuffle w.r.t. Σ_1 and Σ_2 , 208
$\Sigma_1 \underline{\parallel\parallel}_{\Sigma_2}$	fair fS-shuffle w.r.t. Σ_1 and Σ_2 , 208
$\Sigma_1 \underline{\parallel\Gamma}_{\Sigma_2}$	rS-shuffle on Γ w.r.t. Σ_1 and Σ_2 , 209
$\Sigma_1 \underline{\parallel\Gamma}_{\Sigma_2}$	fair rS-shuffle on Γ w.r.t. Σ_1 and Σ_2 , 209
$\ \cdot\ _{i \in [n]}^\Gamma$	n -ary fair S-shuffle on Γ , 227
$\ \cdot\ _{i \in [n]}^\Gamma$	n -ary S-shuffle on Γ , 227
$\underline{\parallel} \cup_{i \in [n]} \Sigma_i$	n -ary fair fS-shuffle w.r.t. $\bigcup_{i \in [n]} \Sigma_i$, 228
$\underline{\parallel} \cup_{i \in [n]} \Sigma_i$	n -ary fS-shuffle w.r.t. $\bigcup_{i \in [n]} \Sigma_i$, 228
$\underline{\parallel\parallel} \cup_{i \in [n]} \Sigma_i$	n -ary fair rS-shuffle on Γ w.r.t. $\bigcup_{i \in [n]} \Sigma_i$, 228
$\underline{\parallel\Gamma} \cup_{i \in [n]} \Sigma_i$	n -ary rS-shuffle on Γ w.r.t. $\bigcup_{i \in [n]} \Sigma_i$, 228

7. Team Automata, I/O Automata, Petri Nets

\mathcal{I}	index set, 233
C_i	component automaton, 233
Σ_i	set of actions (of C_i), 233
\mathcal{S}	composable system, 233
\mathcal{T}	team automaton, 233

Σ	set of actions (of \mathcal{T}), 233
Σ_{ext}	set of external actions (of \mathcal{T}), 233
Σ_{loc}	set of locally-controlled actions (of \mathcal{T}), 233
Θ	arbitrary alphabet disjoint from set Q of states (of \mathcal{T}), 233
\mathcal{S}	compatible system, 237
\mathcal{T}	team I/O automaton, 239
IOCA	$\{\mathbf{B}_C^\Gamma \mid \Gamma \text{ is an alphabet and } C \text{ is a finite input-enabling component automaton with alphabet } \Gamma\}$, 240
IOCA^{alph}	$\{\mathbf{B}_C^{alph} \mid C \text{ is a finite input-enabling component automaton}\}$ (with $alph \in \{inp, out, int, ext, loc\}$), 240
$\Delta_a^v(\mathcal{S})$	complete vector transition space (of a in \mathcal{S}), 245
\underline{a}	vector action a , 245
\mathcal{T}^v	vector team automaton, 245
δ^v	set of labeled vector transitions (of \mathcal{T}^v), 245
$\delta_{\underline{a}}^v$	set of vector \underline{a} -transitions (of \mathcal{T}^v), 245
$SUB_J(\mathcal{T}^v)$	the subteam of \mathcal{T}^v determined by J , 246
\mathcal{T}_F^v	the flattened version (of \mathcal{T}^v), 247
$tFree(\mathcal{T}^v)$	set of truly <i>free</i> actions (of \mathcal{T}^v), 250
$tAI(\mathcal{T}^v)$	set of truly <i>ai</i> actions (of \mathcal{T}^v), 250
$tSI(\mathcal{T}^v)$	set of truly <i>si</i> actions (of \mathcal{T}^v), 250
Λ	empty word vector, 252
$\text{tot}(\{\Delta_j \mid j \in J\})$	total vector alphabet (over $\{\Delta_j \mid j \in J\}$), 252
Δ^u	subset of uniform vector letters of vector alphabet Δ , 252
$v \circ w$	component - wise concatenation (of two n - dimensional vector letters v and w), 252
coll	collapse of a sequence of vector letters into a word vector, 252
$\text{und}(\mathcal{T}^v)$	underlying vector automaton (of \mathcal{T}^v), 253
$\mathbf{V}_{\mathcal{T}^v}$	finitary vector behavior (of \mathcal{T}^v), 253
$\mathbf{V}_{\mathcal{T}^v}^\omega$	infinitary vector behavior (of \mathcal{T}^v), 253
$\mathbf{V}_{\mathcal{T}^v}^\infty$	vector behavior (of \mathcal{T}^v), 253
\mathcal{N}	n -VLITN, 254
P	finite set of places (of \mathcal{N}), 254
T	finite set of events (of \mathcal{N}), 254
O	finite set of n integers, called tokens (of \mathcal{N}), 254
F	flow function (of \mathcal{N}), 254
V	vector alphabet of vector labels (of \mathcal{N}), 255
ℓ	event labeling homomorphism (of \mathcal{N}), 255
$\text{use}(t)$	set of tokens used (by event t), 255
$\mathbf{M}_{\mathcal{N}}$	set of all markings of \mathcal{N} , 255
$\mu[t]_{\mathcal{N}}$	enabled (an event t of \mathcal{N} at a marking μ of \mathcal{N}), 256
$\mu[t]_{\mathcal{N}}\nu$	fires (an event t of \mathcal{N} from a marking μ of \mathcal{N} to a marking ν of \mathcal{N}), 256

$\mu_0[t_1 t_2 \cdots t_m]_{\mathcal{N}}$	firing sequence (of events t_1, t_2, \dots, t_m) of \mathcal{N} starting from μ_0 , 256
$\mu_0[t_1 t_2 \cdots t_m]_{\mathcal{N}} \mu_m$	firing sequence (of events t_1, t_2, \dots, t_m) of \mathcal{N} starting from μ_0 and leading to μ_m , 256
$\mu_0[t_1 t_2 \cdots]_{\mathcal{N}}$	infinite firing sequence (of events t_1, t_2, \dots) of \mathcal{N} starting from μ_0 , 256
\mathcal{K}	n -ITNC, 256
$\text{und}(\mathcal{K})$	underlying n -VLITN (of \mathcal{K}), 256
\mathcal{M}_0	set of initial markings (of \mathcal{K}), 256
\mathcal{M}_f	set of final markings (of \mathcal{K}), 256
$\mathbf{FS}_{\mathcal{K}}$	set of all firing sequences (of \mathcal{K}), 257
$\mathbf{M}_{\mathcal{K}}$	the set of all reachable markings (of \mathcal{K}), 257
$\mathbf{B}_{\mathcal{K}}$	behavior of \mathcal{K} , 257
$\mathbf{V}_{\mathcal{K}}$	vector behavior of \mathcal{K} , 257
carrier (\underline{a})	carrier (of \underline{a}), 260
$PN(\mathcal{T}^v)$	ITNC obtained from \mathcal{T}^v , 261
$SUB_J(\mathcal{K})$	the subnet (of \mathcal{K}) determined by J , 270

8. Applying Team Automata

\mathcal{I}	index set, 278
\mathcal{C}_i	component automaton, 278
$\Sigma_{i,\text{ext}}$	set of external actions (of \mathcal{C}_i), 278
\mathcal{S}	composable system, 278
\mathcal{T}	team automaton, 278
Σ	set of actions (of \mathcal{T}), 278
Σ_{ext}	set of external actions (of \mathcal{T}), 278
\mathcal{C}_H^{Δ}	the Δ -hiding version (of \mathcal{C}), 278
Σ_{com}	set of communicating actions (in \mathcal{S}), 279
$\boxed{\mathcal{T}}$	(communication) closed version (of \mathcal{T}), 279
\mathcal{C}_N^h	h -renamed version (of \mathcal{C}), 280

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