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Structured parallel programming for Monte Carlo Tree Search

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FOR
Monte Carlo Tree Search

S. Ali Mirsoleimani

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FOR

Monte Carlo Tree Search

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I would like to dedicate this thesis to my wife Elahe and to my parents,
for all of their love and support.

In loving memory of my grandfathers,
Bahram and Abolghasem

Preface

The thesis is part of a bigger project, the HEPGAME (High Energy Physics Game). The project started in 2011 when Jos Vermaseren developed the first ideas on improving FORM at Nikhef, Amsterdam. In 2012 he submitted an ERC advanced research grant together with Tilburg University. It was accepted on 12/12/2012. Half a year later in July 2013, the program started. The main objective for HEPGAME was the utilization of AI solutions, particularly by using MCTS for simplification of HEP calculations. One of the issues is solving mathematical expressions of interest with millions of terms. Up to 2011, these calculations were executed with the FORM program, which is software for symbolic manipulation. These calculations are computationally intensive and take a large amount of time. Hence, the FORM program was parallelized to solve large equations in a reasonable amount of time. Therefore, any new algorithm, for instance, the ones based on MCTS, should also be parallelized. Here our research comes in. It is dedicated to parallelization of MCTS on multi-core and many-core processors. The research was ambitious and challenging. Therefore, we divided the research area into three main parts: (1) the evaluation of current methods for parallelization of MCTS, (2) addressing the shortcomings in these methods, and (3) providing new ways of parallelization for MCTS. In the first part, we investigated the current methods and evaluated them in terms of performance and scalability on both multi-core and manycore processors. In the second part, we examined how we can solve the actual shortcomings in the existing parallelization methods for MCTS. The third part was dedicated to finding new ideas, methods, and ways beyond the existing ones to parallelize MCTS.

Sayyed Ali Mirsoleimani, Leiden, July 2019

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List of Abbreviations

3PMCTS	Pipeline Pattern for Parallel MCTS.
FIFO	First In, First Out.
FMA	Fused Multiply Add.
GFLOPS	Giga Floating Point Operations per Second.
GIPS	Giga Integers per Second.
GSCPM	Grain Size Controlled Parallel MCTS.
HEP	High Energy Physics.
HEPGAME	High Energy Physics Game.
ILD	Iteration-Level Dependency.
ILP	Iteration-Level Parallelism.
ILT	Iteration-Level Task.
ISA	Instruction Set Architecture.
MC	Memory Controller.
MCTS	Monte Carlo Tree Search.
MIC	Many Integrated Core.
NUMA	Non Uniform Memory Access.

OLD	Operation-Level Dependency.
OLP	Operation-Level Parallelism.
OLT	Operation-Level Task.
PPS	Playouts per Second.
PS	Problem Statement.
PW	Percentage of Wins.
RNG	Random Number Generation.
RQ	Research Question.
SMT	Simultaneous Multithreading.
TBB	Threading Building Blocks.
TD	Tag Directories.
TPFIFO	Thread Pool with FIFO scheduling.
UCB	Upper Confidence Bound.
UCT	Upper Confidence Bounds for Trees.
UMA	Uniform Memory Access.
VPU s	Vector Processing Units.